

S. Das on Category Theory II

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Def A monad on a category \mathcal{C} is a functor $T: \mathcal{C} \rightarrow \mathcal{C}$ with natural transformations $\eta: 1_{\mathcal{C}} \Rightarrow T$ and $\mu: T^2 \Rightarrow T$ such that the following commute

$$\begin{array}{ccc} T \xrightarrow{\eta} T & \xrightarrow{T\eta} & T \\ \mu T \downarrow & & \downarrow \mu \\ T^2 & \xrightarrow{\mu} & T \end{array} \quad \begin{array}{ccc} T \xrightarrow{\eta T} & T^2 & \\ T\eta \downarrow & \downarrow \eta & \\ T^2 & \xrightarrow{\mu} & T \end{array}$$

Examples 1) $\mathcal{C} = \text{Set}$ $T: A \mapsto \mathcal{P}(A)$ power set
 $\eta: a \mapsto \{a\}$
 $\mu_A: \mathcal{P}(\mathcal{P}(A)) \rightarrow \mathcal{P}(A)$
 set of sets \mapsto union of sets

2) Given an adjunction $F: \mathcal{C} \rightleftarrows \mathcal{D}: G$
 let $T = GF$, $\eta = \text{counit}$, $\mu = G \varepsilon F$ where ε is the unit of $F \dashv G$.

A T-algebra for (T, η, μ) is a pair (X, h) where $X \in \mathcal{C}$ and $h: T(X) \rightarrow X$ with

$$\begin{array}{ccc} T(T(X)) & \xrightarrow{T(h)} & T(X) \\ \mu_X \downarrow & & \downarrow h \\ T(X) & \xrightarrow{\mu} & X \end{array} \quad \begin{array}{ccc} & \xrightarrow{\eta_X} & T(X) \\ X & \xrightarrow{1} & \downarrow h \\ & & X \end{array}$$

The T-algebras form a category \mathcal{C}^T underlying set of

Example 1) Groups $\mathcal{C} = \text{Set}$ $T(X) = \text{free gp on } X$

$$\begin{array}{ccc} \eta_X: X \rightarrow T(X) & & T(T(X)) \rightarrow T(X) \\ x \mapsto [x] & & \text{homomorphism} \end{array}$$

A suitable map $h: T(X) \rightarrow X$ defines a gp structure on X , so T-algebras are gps.

2) Group actions $\mathcal{C} = \text{Set}$, G a gp
 $T(X) = G \times X$ $\mu: (g_1, g_2, x) \mapsto (g_1 g_2, x)$
 $\eta(x) = (e, x)$

A T -algebra is a G -set.

3) R -modules $\mathcal{C} = \text{Ab}$, $T(X) = R \otimes X$

Thm (Eilenberg - Moore construction)

Let (T, η, μ) be a monad on \mathcal{C} . Then the forgetful functor $U: \mathcal{C}^T \rightarrow \mathcal{C}$ has a left adjoint $F: \mathcal{C} \rightarrow \mathcal{C}^T$ sending X to the "free T -algebra" on X .

$X \mapsto (T(X), \mu_X)$. The monad for $F \dashv U$ is (T, η, μ) . Given another adjunction $F': \mathcal{C} \rightleftarrows \mathcal{D}: G'$ with monad (T', η', μ') $F': \mathcal{C} \rightarrow \mathcal{C}^T$ with $F = KF'$ and $G' = UK$.

Limits + colimits Let $X: \mathcal{J} \rightarrow \mathcal{C}$ be a functor to \mathcal{C} from a small cat \mathcal{J} . $j \mapsto X_j$ for $j \in \mathcal{J}$ and a morphism $j \xrightarrow{f} j'$ in \mathcal{J} gives $X_j \xrightarrow{X_f} X_{j'}$. Then $\text{colim}_{\mathcal{J}} X$ is an object W in \mathcal{C} with maps $w_j: X_j \rightarrow W$ with $w_j f = w_{j'}$ and for any other W' with such properties, $\exists! W \rightarrow W'$. Limits are dually defined.

Let \mathcal{C}^J denote the category of functors $J \rightarrow \mathcal{C}$
 We have constant functor $\Delta: \mathcal{C} \rightarrow \mathcal{C}^J$
 then $\text{colim}: \mathcal{C}^J \rightarrow \mathcal{C}$ is left adj of Δ
 adjoints of Δ . lim is right adjoint of
 $\mathcal{C} \rightarrow \mathcal{C}^{J^{op}}$

Examples 1) $J = \text{empty category}$

$\text{lim}_J = \phi = \text{initial object}$

$\text{colim}_J = * = \text{terminal object}$

2) $J = 1 \text{ object category for gp } G$

A functor $J \rightarrow \text{Set or Top}$ defines
 a G -action on a set or space X .

$\text{lim}_J X = X^G = \text{fixed point space}$

$\text{colim}_J X = X_G = \text{orbit space}$

3) $J = \bullet \rightrightarrows \bullet$. A functor $J \rightarrow \mathcal{C}$ is

diagram $A \rightrightarrows B$. $\text{colim}_J = \text{coequalizer}$
 and $\text{lim}_J = \text{equalizer}$.

$\text{lim}_J \rightarrow A \rightrightarrows B \rightarrow \text{colim}_J$.

Every (co)limit is a (co)equalizer.

Given $X: J \rightarrow \mathcal{C}$ define

$$\prod_{j \in \text{Ob } J} X_j \xrightarrow[\theta]{\beta} \prod_{\substack{u: j \rightarrow k \\ \text{in } J}} X_k$$

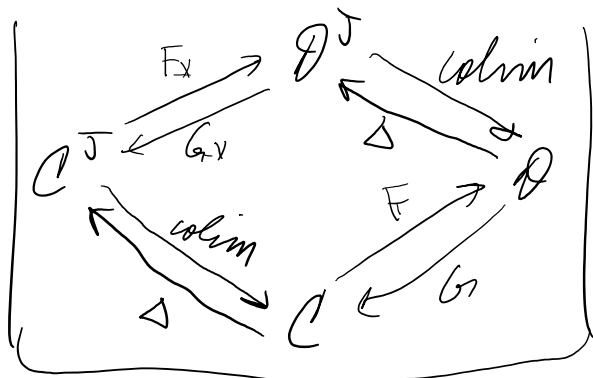
by $\beta_u \beta = \beta_k$ and $\beta_u \theta = X_u \beta_j$. The equalizer
 is $\text{lim}_J X$. Similarly for colimits.

Left (right) adjoints preserve colimits (limits).

Prop Given adjoint functors $F: \mathcal{C} \rightleftarrows \mathcal{D}: G$
 and small cat J , we have $F_x: \mathcal{C}^J \rightleftarrows \mathcal{D}^J: G_x$
 and the map $\text{colim}_J F_x X \xrightarrow{\cong} F \text{colim}_J X$ in \mathcal{D}

Proof For $X \in \mathcal{C}^J$ and $Y \in \mathcal{D}^J$
 $\mathcal{D}^J(F_x X, Y) = \mathcal{D}^J(FX, Y)$ where $J \xrightarrow{X} \mathcal{C} \xrightarrow{F} \mathcal{D}$
 $\mathcal{C}^J(X, G_x Y) = \mathcal{C}^J(X, GY)$ $J \xrightarrow{Y} \mathcal{D} \xrightarrow{G} \mathcal{C}$
 Set $\theta \in \mathcal{D}^J(FX, Y)$, as natural terms $FX \Rightarrow Y$
 $\theta \in \mathcal{C}^J(X, GY)$ " $X \Rightarrow GY$
 $\theta_j: X_j \rightarrow GY_j \iff \theta'_j: FX_j \rightarrow Y_j$ since $F \dashv G$.

Claim this diagram commutes



For $X \in \mathcal{C}^J$ and $Y \in \mathcal{D}^J$
 $\mathcal{C}^J(X, G_x \Delta Y) \cong \mathcal{D}^J(F_x X, \Delta Y)$
 $\cong \mathcal{C}(\text{colim}_J F_x X, Y)$
 Hence.
 $F_x \text{colim}_J \dashv \Delta G$
 so $F_x \text{colim}_J$ and $\text{colim}_J F_x$

are adjoint to the same functors and are therefore equal QED

A reflexive coequalizer is the colimit of



A reflexive equalizer is its limit.

Such a diagram is a functor from $J = \begin{matrix} \bullet & \rightrightarrows & \bullet \\ \lrcorner & & \lrcorner \end{matrix}$

It has subcats $\text{End}_A \mathcal{B}$ and $\tilde{J} = \begin{matrix} \bullet & \rightrightarrows & \bullet \\ \lrcorner & & \lrcorner \end{matrix}$

Then $\text{colim}_J = \text{colim}_{\tilde{J}} = \text{colim}_{\text{End}_A}$